

# Microeconomics IV (Game Theory)

## Lecture 5 - Some Refinements of Equilibrium in Strategic Form

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2026

# Refinements of Nash Equilibrium

# Introduction

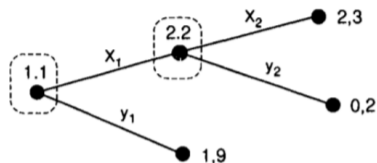
- Any extensive-form game can be represented in strategic form.
- There are multiple representations:
  - **Normal representation**
  - **Reduced normal representation**
  - **Multiagent representation**
- These forms raised the hope that all principles of game-theoretic analysis could be applied in the simpler strategic form.

# Strategic vs Extensive Form

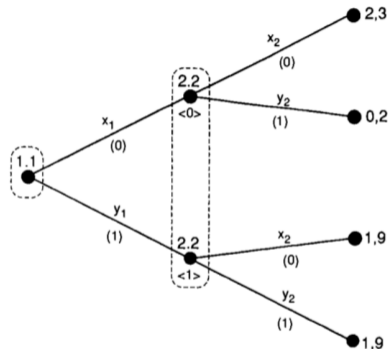
- However, this hope does **not always hold**.
- It is easy to construct extensive-form games that:
  - Are identical in strategic form (under any representation),
  - But yield **different sets of sequential equilibria**.

# Example 1: Same Strategic Form, Different Equilibria

## Game 1



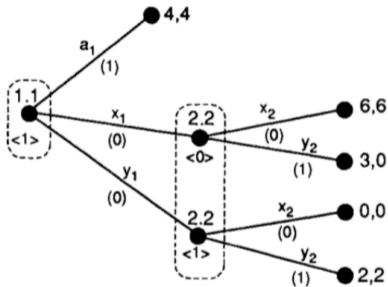
## Game 2



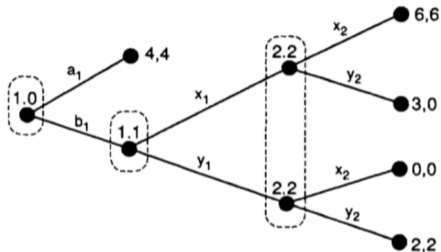
- Both games have **identical strategic representations**.
- But  $([y_1], [y_2])$  is a **sequential equilibrium** in Game 2, not in Game 1.
- Extensive form matters for belief formation and rationality.

## Example 2: Games 3 and 4

### Game 3



### Game 4



- **Game 3** and **Game 4** have the *same reduced normal form* (up to relabeling).
- But: Game 3 admits a sequential equilibrium with Player 1 choosing  $a_1$ .
- In Game 4, no such sequential or subgame-perfect equilibrium exists.

## Why the Difference?

- In **Game 4**, Player 2 observes that Player 1 did *not* choose  $a_1$ .
- So Player 2 must infer Player 1 chose between  $x_1$  and  $y_1$  *intentionally*.
- Rationality  $\Rightarrow$  Player 1 would choose the better move between  $x_1$  and  $y_1 \Rightarrow$  surely  $x_1$ .
- So Player 2 chooses  $x_2$ .
- This logic prevents  $a_1$  from being part of any equilibrium.

# Why Refinements Are Needed

- We saw that:
  - Strategic-form equivalence can mask **differences in sequential behavior**.
  - Some sequential equilibria include **weakly dominated strategies**.
  - Some games have **no reasonable sequential equilibrium**.
- To resolve these issues, we need **refinements of Nash equilibrium** that:
  - Apply in strategic form.
  - Reflect **sequential rationality** from the extensive form.

# Desirable Properties of a Refined Equilibrium

- A refined equilibrium should:
  - Always **exist** for any finite game.
  - Exclude outcomes involving **dominated strategies**.
  - Eliminate equilibria that **cannot be justified by any belief system**.
  - Preserve equivalence across:
    - Normal form,
    - Reduced normal form,
    - Multiagent representations.

## Core Idea

### **i** Note

A good refinement of Nash equilibrium should give the same solution across all strategic-form representations of the same extensive-form game.

- This ensures that solution concepts are **representation-invariant**.

# Perfect Equilibrium

# Perfect Equilibrium

- One such refinement: **Perfect Equilibrium** (Selten, 1975).
- Key features:
  - Based on **trembles** (small mistakes with positive probability).
  - Eliminates equilibria that rely on **non-credible threats**.
  - Preserves rationality across strategic and extensive forms.
- Especially robust when applied through the **multiagent representation**.

# Perfect Equilibrium

## Perfect Equilibrium

Let  $\Gamma = (N, (C_i)_{i \in N}, (u_i)_{i \in N})$  be a finite game in strategic form. A mixed-strategy profile  $\sigma \in \times_{i \in N} \Delta(C_i)$  is a **perfect equilibrium** if there exists a sequence  $(\hat{\sigma}^k)_{k=1}^{\infty}$  such that:

$$\bullet \hat{\sigma}^k \in \times_{i \in N} \Delta^0(C_i), \quad \forall k \in \{1, 2, 3, \dots\} \quad (1)$$

$$\bullet \lim_{k \rightarrow \infty} \hat{\sigma}_i^k(c_i) = \sigma_i(c_i), \quad \forall i \in N, \forall c_i \in C_i \quad (2)$$

$$\bullet \sigma_i \in \arg \max_{\tau_i \in \Delta(C_i)} u_i(\hat{\sigma}_{-i}^k, \tau_i), \quad \forall i \in N \quad (3)$$

- For any finite set  $Z$ ,  $\Delta^0(Z) = \{q \in \Delta(Z) \mid q(z) > 0, \forall z \in Z\}$  – the set of probability distributions assigning strictly positive probability to every element in  $Z$ .

# Interpretation

- (1): Each  $\hat{\sigma}^k$  assigns **positive probability to every pure strategy**.
- (2): The sequence converges to the equilibrium  $\sigma$ .
- (3): Each  $\sigma_i$  is a **best response** to  $\hat{\sigma}_{-i}^k$ .

## **i** Note

Perfect Equilibrium refines Nash by ensuring robustness to small trembles in opponents' strategies.

# Perfect Equilibrium Implies Nash

- Every perfect equilibrium is also a **Nash equilibrium**.
- Why? Because:
  - The **best-response correspondence** is upper hemicontinuous.
  - Limits of perturbed equilibria still satisfy the best-response condition.
- Formally:

$$\sigma_i \in \arg \max_{\tau_i \in \Delta(C_i)} u_i(\sigma_{-i}, \tau_i), \quad \forall i \in N$$

## **i** Note

In a perfect equilibrium, strategies remain best responses even when every pure strategy is played with positive probability in the perturbation sequence.

## Connection to Sequential Equilibrium

- Both **perfect** and **sequential** equilibrium test  $\sigma$  against **perturbed strategies**.
- In sequential equilibrium:
  - Perturbations help generate beliefs at zero-probability information sets.
  - Beliefs are used in the optimality condition.
- In perfect equilibrium:
  - Perturbations are directly used in the **payoff comparison** (condition 3).
  - No need to construct beliefs – just evaluate payoffs from trembles.

### **i** Note

In the multiagent representation, condition (3) is stronger than the sequential rationality condition. Therefore,  $\Rightarrow$  Every perfect equilibrium of the multiagent game  $\Gamma^e$  is a sequential equilibrium of  $\Gamma^e$ .

## Recall Notation – Game, Players, Nodes

Symbol	Meaning
$\Gamma^e$	Extensive-form game with <b>perfect recall</b>
$i, s \in S_i$	Player $i$ and one of his information states $s$
$D_s$	Set of moves available at state $s$
$Y_s$	Set of <b>decision nodes</b> in information state $s$
$X(s)$	Set of terminal nodes that do <b>not</b> follow any node in $Y_s$ – the <b>bypass set</b> (plays that never reach $s$ )
$w_i(x)$	Payoff to player $i$ at terminal node $x$

- $X(s)$  is the one piece worth pausing on: it collects exactly the plays where  $i$ 's move at  $s$  is **irrelevant**, because the game never visits  $s$ .

# Recall Notation – Strategies, Beliefs, Payoffs

Symbol	Meaning
$\hat{\sigma}^k$	A sequence of <b>fully mixed</b> (strictly positive) profiles supporting $\sigma$ as a perfect equilibrium of the multiagent rep.
$\rho_s \in \Delta(D_s)$	An alternative local move distribution at state $s$
$\bar{P}(y   \sigma)$	Probability that node $y$ is reached under profile $\sigma$
$U_i(\sigma   y)$	Expected <b>continuation</b> payoff to $i$ from node $y$ onward
$v_s$	Payoff of player $i$ 's <b>agent at state</b> $s$ in the multiagent rep – an <b>ex ante</b> (whole-tree) expectation, <i>not</i> a conditional one
$\hat{\pi}_s^k, \pi_s$	Beliefs from $\hat{\sigma}^k$ by Bayes; $\pi_s = \lim_k \hat{\pi}_s^k$

## Theorem: Perfect $\Rightarrow$ Sequential

### Theorem 1

Suppose  $\Gamma^e$  is an extensive-form game with **perfect recall** and  $\sigma$  is a **perfect equilibrium** of the **multiagent representation** of  $\Gamma^e$ . Then there exists a belief vector  $\pi$  such that  $(\sigma, \pi)$  is a **sequential equilibrium** of  $\Gamma^e$ .

- Links perfect equilibrium (a strategic-form refinement) to sequential equilibrium (an extensive-form refinement).
- Both are built from the **same trembling-hand limit** – perfection's node-by-node best response is exactly the raw material sequential rationality needs.

## Proof Sketch (1/3) – Beliefs from Trembles

### **i** Note

**Goal:** From a perfect equilibrium  $\sigma$  of the multiagent rep, build beliefs  $\pi$  so that  $(\sigma, \pi)$  is a sequential equilibrium.

- Perfection gives a sequence of **fully mixed** profiles  $\hat{\sigma}^k \rightarrow \sigma$ .
- For each  $k$  and each node  $y \in Y_s$ , define beliefs by Bayes' rule:

$$\hat{\pi}_s^k(y) = \frac{\bar{P}(y | \hat{\sigma}^k)}{\sum_{z \in Y_s} \bar{P}(z | \hat{\sigma}^k)}$$

- Since  $\hat{\sigma}^k$  is strictly positive,  $\bar{P}(y | \hat{\sigma}^k) > 0$  for every  $y$ , so the denominator is **positive** and the ratio is well-defined.
- Set  $\pi_s(y) = \lim_{k \rightarrow \infty} \hat{\pi}_s^k(y)$ . By construction  $\pi$  is **fully consistent** with  $\sigma$ .

## Proof Sketch (2/3) – Splitting the Agent's Payoff

- The agent- $s$  payoff against  $\hat{\sigma}^k$ , when playing  $\rho_s$ , splits into plays **through**  $s$  and plays **bypassing**  $s$ :

$$v_s(\hat{\sigma}_{-i,s}^k, \rho_s) = \sum_{y \in Y_s} \bar{P}(y | \hat{\sigma}^k) U_i(\hat{\sigma}_{-i,s}^k, \rho_s | y) + \sum_{x \in X(s)} \bar{P}(x | \hat{\sigma}^k) w_i(x)$$

- Key fact:**  $\bar{P}(y | \hat{\sigma}_{-i,s}^k, \rho_s) = \bar{P}(y | \hat{\sigma}^k)$  – reaching  $y \in Y_s$  depends only on agents who move **before**  $s$ , not on  $\rho_s$ .
- The bypass term  $\sum_{x \in X(s)} \bar{P}(x | \hat{\sigma}^k) w_i(x)$  is **constant in**  $\rho_s$  – those plays never reach  $s$ .

## Proof Sketch (2/3, cont.) – Factoring out Beliefs

- Multiply and divide the first sum by the reach-probability of  $s$ :

$$v_s = \underbrace{\left( \sum_{y \in Y_s} \hat{\pi}_s^k(y) U_i(\hat{\sigma}_{-i,s}^k, \rho_s | y) \right)}_{\text{belief-weighted continuation}} \underbrace{\left( \sum_{z \in Y_s} \bar{P}(z | \hat{\sigma}^k) \right)}_{\text{reach prob. of } s > 0} + \underbrace{\sum_{x \in X(s)} \bar{P}(x | \hat{\sigma}^k) w_i(x)}_{\text{const. in } \rho_s}$$

- So  $v_s$  is a **strictly increasing affine transformation** of the conditional, belief-weighted objective – with coefficients **independent of**  $\rho_s$ .
- $\rightarrow$  The two objectives have the **same maximizers**.

## Proof Sketch (3/3) – Sequential Rationality

- Perfection means  $\sigma_s$  best-responds in the agent's payoff at every  $k$ :

$$\sigma_s \in \arg \max_{\rho_s \in \Delta(D_s)} v_s(\hat{\sigma}_{-i,s}^k, \rho_s)$$

- By the affine equivalence, equivalently:

$$\sigma_s \in \arg \max_{\rho_s \in \Delta(D_s)} \sum_{y \in Y_s} \hat{\pi}_s^k(y) U_i(\hat{\sigma}_{-i,s}^k, \rho_s \mid y)$$

- Let  $k \rightarrow \infty$ . By **upper-hemicontinuity** of the best-response correspondence (limits of maximizers maximize the limit):

$$\sigma_s \in \arg \max_{\rho_s \in \Delta(D_s)} \sum_{y \in Y_s} \pi_s(y) U_i(\sigma_{-i,s}, \rho_s \mid y)$$

- This is exactly the sequential-rationality condition. With  $\pi$  fully consistent (Step 1),  $(\sigma, \pi)$  is a sequential equilibrium. ■

## Selten's Definition (1975)

### **i** Note

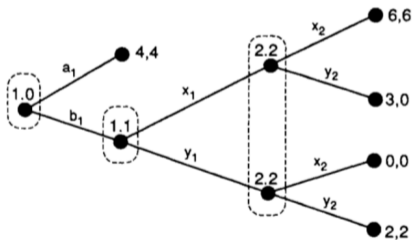
Selten (1975) defined a **perfect equilibrium** of an extensive-form game to be any perfect equilibrium of its **multiagent representation**.

- This **contrasts** with the definition of **Nash** equilibrium:
  - A Nash equilibrium of an extensive-form game is defined as a Nash equilibrium of its multiagent representation,
  - **Only if** the corresponding mixed strategy profile is also a Nash equilibrium in the **normal representation**.
- Perfect equilibrium relies only on the multiagent structure, not the normal form.

## Example 3

- Consider **Game 4** and its normal representation:

Game 4



Normal Representation

$C_1$	$C_2$	
	$x_2$	$y_2$
$a_1x_1$	4,4	4,4
$a_1y_1$	4,4	4,4
$b_1x_1$	6,6	3,0
$b_1y_1$	0,0	2,2

## Perfect Equilibrium in the Normal Form

- The strategy profile  $([a_1x_1], [y_2])$  is a **perfect equilibrium** in the normal representation.
- To see this, construct the perturbed strategy sequence:

$$\hat{\sigma}^k = ((1 - \varepsilon)[a_1x_1] + .1\varepsilon[a_1y_1] + .1\varepsilon[b_1x_1] + .8\varepsilon[b_1y_1], \quad \varepsilon[x_2] + (1 - \varepsilon)[y_2])$$

with  $\varepsilon = 1/(k + 2)$ .

## Numerical Justification

- When  $\varepsilon \leq \frac{1}{3}$ :

- Player 1:

$$\text{Payoff from } a_1x_1 = 4 \geq 6\varepsilon + 3(1 - \varepsilon)$$

So  $[a_1x_1]$  is a best response to  $\varepsilon[x_2] + (1 - \varepsilon)[y_2]$ .

- Player 2:

- If  $b_1y_1$  has higher probability (more than 3 times larger) than  $b_1x_1$ , then  $y_2$  is optimal.

### **i** Note

Hence,  $(\hat{\sigma}^k)$  satisfies (1)–(3), so  $([a_1x_1], [y_2])$  is a perfect equilibrium in the normal form.

## Multiagent Representation Breakdown

- In the multiagent form, equilibrium must hold **at each node**, for each **agent**.
- Suppose Player 1's agent at node 1.0 plays  $b_1$  with positive probability.
- Then at node 1.1, the unique best response must be  $x_1$ :

$$\lim_{k \rightarrow \infty} \hat{\sigma}_{1.1}^k(x_1) = 1, \quad \hat{\sigma}_{1.1}^k(y_1) = 0$$

- This implies:
  - $b_1 x_1$  occurs with high probability,
  - So Player 2 responds with  $x_2$ , not  $y_2$ .

## Consequences for Perfect Equilibrium

- If  $x_1$  and  $x_2$  are highly likely:
  - The best response at 1.0 is  $b_1$  (since  $6 > 4$ ),
  - So Player 1 plays  $b_1$  with high probability.
- Therefore, any perfect equilibrium in the **multiagent representation** must satisfy:

$$([b_1], [x_1], [x_2])$$

### **i** Note

⇒ The strategy profile  $([a_1 x_1], [y_2])$  is a perfect equilibrium in the normal form, but **not** in the multiagent (or extensive-form) representation.

# Existence of Perfect and Sequential Equilibria

## Theorem 2

For any finite game in strategic form, there exists at least one perfect equilibrium.

- Theorems **1** and **2** immediately imply the existence theorem for sequential equilibria.

## Example 4 - Perfect Equilibrium in Normal Games

- Consider a 3-player game with:
  - Player 1 strategy set:  $S_1 = \{U, D\}$
  - Player 2 strategy set:  $S_2 = \{L, R\}$
  - Player 3 strategy set:  $S_3 = \{X, Y\}$

**Player 3 plays  $X$**

	$L$	$R$
$U$	$(1, 1, 1)$	$(1, 0, 1)$
$D$	$(1, 1, 1)$	$(0, 0, 1)$

**Player 3 plays  $Y$**

	$L$	$R$
$U$	$(1, 1, 0)$	$(0, 0, 0)$
$D$	$(0, 1, 0)$	$(1, 0, 0)$

## Example 4 - Perfect Equilibrium in Normal Games

- Perform iterated deletion of weakly dominated strategies. What are the strategies that survive this process?
- **Player 2:**
  - Strategy  $L$  is strictly dominant over  $R$
  - Keep  $L$ , delete  $R$
- **Player 3:**
  - Strategy  $X$  is strictly dominant over  $Y$
  - Keep  $X$ , delete  $Y$
- **Player 1:**
  - No strict or weak domination between  $U$  and  $D$
  - No strategies deleted for player 1
- **Surviving Strategies:**  $\{\{U, D\}, \{L\}, \{X\}\}$

## Example 4 - Perfect Equilibrium in Normal Games

- Find all Nash equilibria (both pure and mixed). Does any of the Nash equilibrium involve weakly dominated strategies?
- **Player 2** plays strictly dominant strategy  $L$
- **Player 3** plays strictly dominant strategy  $X$
- Given  $(L, X)$ , Player 1 is **indifferent** between  $U$  and  $D$
- So any mixed strategy of Player 1 between  $U$  and  $D$  is a best response
- **Nash Equilibria:**

$$(pU + (1 - p)D, L, X), \quad \text{for any } p \in [0, 1]$$

- **None of these NE involve weakly dominated strategies**

## Example 4 - Perfect Equilibrium in Normal Games

- Focus on the pure strategy Nash equilibria. Are they trembling hand perfect?
- The game has two pure strategy Nash equilibria:

$$(U, L, X) \quad \text{and} \quad (D, L, X)$$

- We will analyze whether these are **trembling-hand perfect** (THP) equilibria (alternative name).

## Example 4 - Perfect Equilibrium in Normal Games

- Let the perturbed (completely mixed) strategy profile be:

$$\sigma_1(\varepsilon_1) = \varepsilon_1 U + (1 - \varepsilon_1) D$$

$$\sigma_2(\varepsilon_2) = (1 - \varepsilon_2) L + \varepsilon_2 R$$

$$\sigma_3(\varepsilon_3) = (1 - \varepsilon_3) X + \varepsilon_3 Y$$

- As  $\varepsilon \rightarrow (0, 0, 0)$ , the profile  $\sigma(\varepsilon) \rightarrow (D, L, X)$

## Example 4 - Perfect Equilibrium in Normal Games

- Checking if  $(D, L, X)$  is THP
- Compute expected payoffs for Player 1:

$$u_1(U, \sigma_2, \sigma_3) = 1 - \varepsilon_2\varepsilon_3 \quad (1)$$

$$u_1(D, \sigma_2, \sigma_3) = 1 - \varepsilon_2 - \varepsilon_3 + 2\varepsilon_2\varepsilon_3 \quad (2)$$

- For  $D$  to be a best response:

$$1 - \varepsilon_2 - \varepsilon_3 + 2\varepsilon_2\varepsilon_3 \geq 1 - \varepsilon_2\varepsilon_3 \Rightarrow 3\varepsilon_2\varepsilon_3 \geq \varepsilon_2 + \varepsilon_3$$

- But as  $\varepsilon_i \rightarrow 0$  for  $i = 2, 3$ , this inequality is eventually violated.
- For example, if  $\varepsilon_2 \in (0, \frac{1}{3})$ , the inequality fails regardless of  $\varepsilon_3$ .
- $\Rightarrow D$  cannot be a best response to the sequence of totally mixed strategies of players 2 and 3.

## Example 4 - Perfect Equilibrium in Normal Games

- Is  $(U, L, X)$  THP?
- From equations (1) and (2), we now check if  $U$  can be a best response.
- We want a sequence of  $(\varepsilon_2, \varepsilon_3) \rightarrow 0$  such that  $U$  is better than  $D$ .
- Condition for  $U$  being a best response:

$$3\varepsilon_2\varepsilon_3 \leq \varepsilon_2 + \varepsilon_3 \quad \text{or} \quad 3\varepsilon_2 \leq 1 + \frac{\varepsilon_2}{\varepsilon_3}$$

- It suffices to let  $\varepsilon_2$  be any sequence:
  - bounded above by  $\frac{1}{3}$ , and
  - converging to 0
- Let  $\varepsilon_3$  be any sequence converging to 0
- Hence,  $U$  can be a best response in the limit, so  $(U, L, X)$  is **trembling-hand perfect**.

## Implication: THP in Multi-Player Games

- In 2-player games:  
THP  $\Leftrightarrow$  no weakly dominated strategies in NE
- In 3+ player games:  
**THP can rule out NE that use only undominated strategies**

### **i** Note

This game shows that some pure-strategy Nash equilibria – like  $(D, L, X)$  – are **not trembling-hand perfect**, even though they involve no weakly dominated strategies.

## Why Trembling-Hand Perfection Makes Sense

- The two equilibria  $(U, L, X)$  and  $(D, L, X)$  give **identical on-path outcomes** – Player 1 is indifferent.
- But indifference holds only on the knife-edge where Players 2 and 3 play  $L, X$  with probability **exactly** 1.
- **Once we admit that hands shake**, the tie breaks:
  - $U$  wins in the **single-tremble** cells  $(R, X)$  and  $(L, Y)$  – first order,  $\sim \varepsilon_2 + \varepsilon_3$ .
  - $D$  wins only in the **double-tremble** cell  $(R, Y)$  – second order,  $\sim \varepsilon_2\varepsilon_3$ .
- Single mistakes are far likelier than simultaneous ones, so  $U$  is the **safe** response.
- *THP is just this intuition made precise*: don't merely best-respond when everyone plays perfectly – best-respond when everyone *almost* does. The catch: it bounds only the **size** of mistakes, not their **relative** likelihood. Constraining *that* – worse mistakes must be rarer – is the step from perfect to **proper** equilibrium.

# Proper Equilibrium

# Proper Equilibria

- Perfect equilibrium is defined using the **multiagent representation** of extensive-form games.
- However, even games with the **same reduced normal form** (e.g., Games 3 and 4) may be distinguished by perfect equilibrium.
- For example:
  - Both  $([a_1], [y_2])$  and  $([x_1], [x_2])$  are perfect equilibria in the multiagent version of **Game 3**
  - But only  $([b_1], [x_1], [x_2])$  is perfect equilibrium in the multiagent form of **Game 4**

# Motivation for Proper Equilibria

- This sensitivity motivates a refinement that **respects sequential rationality**, even when applied to the normal form.
- We want a solution concept that:
  - Selects equilibria that **can be extended to sequential equilibria**
  - Does **not distinguish games** based on artifacts of representation

## **i** Note

Proper equilibria form a **subset of perfect equilibria**, and always correspond to equilibria that can be extended to **sequential equilibria**.

## $\varepsilon$ -Perfect Equilibrium

- Let  $\Gamma = (N, (C_i)_{i \in N}, (u_i)_{i \in N})$  be a finite normal-form game.
- Let  $\Delta^0(C_i)$  be the set of completely mixed strategies for player  $i$ .
- A strategy profile  $\sigma \in \times_{i \in N} \Delta^0(C_i)$  is an  $\varepsilon$ -perfect equilibrium if:

For all  $i \in N$  and all  $c_i \in C_i$  : if  $c_i \notin \arg \max_{e_i \in C_i} u_i(\sigma_{-i}, [e_i])$ , then  $\sigma_i(c_i) < \varepsilon$

- In an  $\varepsilon$ -**perfect equilibrium**, players are allowed to make small mistakes.

## Limit Characterization of Perfect Equilibrium

- A strategy profile  $\bar{\sigma}$  is a perfect equilibrium if and only if there exists a sequence  $(\varepsilon(k), \hat{\sigma}^k)_{k=1}^{\infty}$  such that:

$$\lim_{k \rightarrow \infty} \varepsilon(k) = 0, \quad \lim_{k \rightarrow \infty} \hat{\sigma}_i^k(c_i) = \bar{\sigma}_i(c_i) \quad \text{for all } i \in N, \text{ and all } c_i \in C_i$$

and each  $\hat{\sigma}^k$  is an  $\varepsilon(k)$ -perfect equilibrium

### **i** Note

Perfect equilibrium is the limit of a sequence of  $\varepsilon$ -perfect equilibria as  $\varepsilon \rightarrow 0$ , where players assign vanishing probability to suboptimal pure strategies.

## $\varepsilon$ -Proper Equilibrium

- Let  $\sigma \in \times_{i \in N} \Delta^0(C_i)$  be a completely mixed strategy profile.
- $\sigma$  is an  $\varepsilon$ -proper equilibrium if:

For all  $i \in N$  and all  $c_i, e_i \in C_i$  : if  $u_i(\sigma_{-i}, [c_i]) < u_i(\sigma_{-i}, [e_i])$ , then  $\sigma_i(c_i) \leq \varepsilon \cdot \sigma_i(e_i)$

- That is, **worse strategies must be played with much smaller probability** than better ones.

## Proper Equilibrium (Limit Characterization)

- A strategy profile  $\bar{\sigma}$  is a **proper equilibrium** if there exists a sequence  $(\varepsilon(k), \sigma^k)$  such that:

$$\lim_{k \rightarrow \infty} \varepsilon(k) = 0, \quad \lim_{k \rightarrow \infty} \sigma_i^k(c_i) = \bar{\sigma}_i(c_i), \quad \forall i \in N, \quad \forall c_i \in C_i$$

- And each  $\sigma^k$  is an  $\varepsilon(k)$ -proper equilibrium.

# Intuition

- Every pure strategy gets positive probability.
- Strategies that are **clear mistakes** (i.e., strictly worse than others) receive **vanishingly small probability**.
- A strictly worse strategy must be played with an asymptotically smaller probability than a strictly better strategy.
- Proper equilibrium is a refinement of perfect equilibrium:
  - **Every proper equilibrium is perfect.**
  - **Not every perfect equilibrium is proper.**
- *Proper equilibrium ensures robustness to mistakes: the worse the mistake, the smaller the chance a rational player makes it.*

# Existence of Proper Equilibrium

## Theorem 3

For any finite game in strategic form, the set of proper equilibria is a nonempty subset of the set of perfect equilibria.

# Proper Equilibria and Sequential Equilibria

## Theorem 4

Suppose that  $\Gamma^e$  is an extensive-form game with perfect recall and  $\tau$  is a proper equilibrium of the normal representation of  $\Gamma^e$ . Then there exists a vector of belief probabilities  $\pi$  and a behavioral strategy profile  $\sigma$  such that  $(\sigma, \pi)$  is a sequential equilibrium of  $\Gamma^e$ , and  $\sigma$  is a behavioral representation of  $\tau$ .

# Interpretation

- Just like **perfect equilibrium** selects **sequential equilibria** in the **multiagent** representation, **proper equilibrium** selects **sequential equilibria** in the **normal-form** representation.
- This bridges the gap between:
  - **Strategic-form refinements** (e.g., proper equilibrium), and
  - **Extensive-form rationality** (e.g., sequential equilibrium)

# Implication

- Every proper equilibrium of the normal form can be interpreted as a **sequentially rational plan** in the extensive form.
- Ensures consistency between:
  - Mistake-aware normal-form refinements, and
  - On-path and off-path beliefs in dynamic games

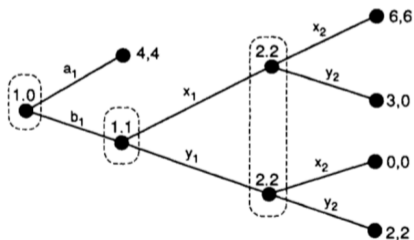
## **i** Note

Proper equilibrium is not just a refinement of Nash – it is also a foundation for constructing **sequential equilibria** from the normal form.

# Back to Game 4

- Consider **Game 4** and its normal representation:

Game 4



Normal Representation

$C_1$	$C_2$	
	$x_2$	$y_2$
$a_1x_1$	4,4	4,4
$a_1y_1$	4,4	4,4
$b_1x_1$	6,6	3,0
$b_1y_1$	0,0	2,2

## Example: $\varepsilon$ -Perfect but Not $\varepsilon$ -Proper

- Consider the strategy profile:

$$\hat{\sigma} = ((1 - \varepsilon)[a_1x_1] + .1\varepsilon[a_1y_1] + .1\varepsilon[b_1x_1] + .8\varepsilon[b_1y_1], \varepsilon[x_2] + (1 - \varepsilon)[y_2])$$

- This profile is  $\varepsilon$ -perfect for any  $\varepsilon < \frac{1}{3}$
- So as  $\varepsilon \rightarrow 0$ ,  $\hat{\sigma}$  converges to  $([a_1x_1], [y_2])$ , a **perfect equilibrium**

## Not $\varepsilon$ -Proper

- Although  $\hat{\sigma}$  is  $\varepsilon$ -perfect, it is **not**  $\varepsilon$ -proper
- Why? Consider the strategies  $b_1x_1$  and  $b_1y_1$  for player 1:
  - $b_1y_1$  is a **worse mistake** than  $b_1x_1$
  - Payoffs:

$$u_1(b_1y_1) = 0\varepsilon + 2(1 - \varepsilon) = 2 - 2\varepsilon$$

$$u_1(b_1x_1) = 6\varepsilon + 3(1 - \varepsilon) = 3 + 3\varepsilon$$

- But in  $\hat{\sigma}$ , the ratio of probabilities is:

$$\frac{\hat{\sigma}_1(b_1y_1)}{\hat{\sigma}_1(b_1x_1)} = \frac{0.8\varepsilon}{0.1\varepsilon} = 8$$

- $\varepsilon$ -properness requires this ratio  $\leq \varepsilon$  – clearly violated

# Unique Proper Equilibrium

- For  $\varepsilon < 1$ , to satisfy  $\varepsilon$ -properness:
  - Player 1 must assign **less probability** to  $b_1 y_1$  than to  $b_1 x_1$
  - So  $x_2$  must be a best response for player 2
- Hence, in any proper equilibrium:
  - Player 2 plays  $x_2$
  - Player 1 plays  $b_1 x_1$
- The unique proper equilibrium of the strategic-form game is:

$$(b_1 x_1, x_2)$$

## Justification by $\varepsilon$ -Proper Equilibria

- This proper equilibrium arises as a limit of  $\varepsilon$ -proper equilibria like:

$$((1 - \varepsilon - .5\varepsilon^2)[b_1x_1] + .5\varepsilon^2[b_1y_1] + .5\varepsilon[a_1x_1] + .5\varepsilon[a_1y_1], (1 - \varepsilon)[x_2] + \varepsilon[y_2])$$

- This sequence respects properness:
  - Mistakes get vanishingly small probability
  - Worse mistakes get **even less** than mild ones

### **i** Note

This proper equilibrium corresponds to the **unique sequential equilibrium** of the extensive-form game 4.

## Finding a Proper Equilibrium

- 1 Find the **Nash equilibria** of the normal-form game.
- 2 For a candidate  $\bar{\sigma}$ , construct completely mixed strategies

$$\sigma^\varepsilon \rightarrow \bar{\sigma}.$$

- 3 Compute and rank the payoffs of all pure strategies against  $\sigma_{-i}^\varepsilon$ .
- 4 Impose

$$u_i(c_i, \sigma_{-i}^\varepsilon) < u_i(e_i, \sigma_{-i}^\varepsilon) \Rightarrow \sigma_i^\varepsilon(c_i) \leq \varepsilon \sigma_i^\varepsilon(e_i).$$

- 5 Verify these inequalities for sufficiently small  $\varepsilon$ .

### **i** Note

If such a sequence exists,  $\bar{\sigma}$  is a **proper equilibrium**.

# Required Reading

- **Myerson (1997).**
  - Chapter 5 – 5.1-5.4